

## Edge eccentric connectivity index of nanothorns

Z. N. Berberler, M. E. Berberler

Faculty of Science, Department of Computer Science, Dokuz Eylul University,  
35160, Izmir, Turkey

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The edge eccentric connectivity index of a graph  $G$  is defined as  $\xi_e^c(G) = \sum_{f \in E(G)} \deg_G(f) ec_G(f)$ , where  $\deg_G(f)$  is the degree of an edge  $f$  and  $ec_G(f)$  is its eccentricity. In this paper, we investigate the edge eccentric connectivity index of a class of nanographs, namely, nanothorns and we present an explicit formula for the edge eccentric connectivity index of nanothorns. The results are applied to compute this eccentricity-related invariant for some important classes of chemical graphs by specializing components of this type of graphs. In addition, an algorithm is designed for calculating the edge eccentric connectivity index of graphs.

**Keywords:** Edge eccentric connectivity index, Eccentric connectivity index, Eccentricity, Zagreb index, Thorny graph.

### INTRODUCTION

Let  $G = (V, E)$  be a connected simple graph with vertex and edge sets  $V(G)$  and  $E(G)$ , respectively. If  $|V(G)| = n$  and  $|E(G)| = q$ , we say that  $G$  is an  $(n, q)$ -graph. The degree of a vertex  $v \in V(G)$ , denoted by  $\deg_G(v)$ , is the number of edges incident to  $v$ . A vertex of degree one is called a *pendent vertex*. The edge incident to a pendent vertex is called a *pendent edge*. The distance  $d_G(u, v)$  between two vertices  $u$  and  $v$  in  $G$  is the length of the shortest path between them. The eccentricity of a vertex  $u \in V(G)$ , denoted by  $ec_G(u)$ , is the largest distance between  $u$  and any other vertex  $v$  of  $G$ . Let  $f = uv$  be an edge in  $E(G)$ . Then the degree of an edge  $f$ , denoted as  $\deg_G(f)$ , is defined to be  $\deg_G(f) = \deg_G(u) + \deg_G(v) - 2$ . For two edges  $f_1 = u_1v_1$ ,  $f_2 = u_2v_2$  in  $E(G)$ , the distance between two edges  $f_1$  and  $f_2$ , denoted as  $d_G(f_1, f_2)$ , is defined to be  $d_G(f_1, f_2) = \min\{d_G(u_1, u_2), d_G(u_1, v_2), d_G(v_1, u_2), d_G(v_1, v_2)\}$ . The eccentricity of an edge  $f$ , denoted by

$ec_G(f)$ , is defined as

$ec_G(f) = \max\{d_G(f, e) \mid e \in E(G)\}$ . The corona product  $G \circ H$  of two graphs  $G$  and  $H$  is defined as the graph obtained by taking one copy of  $G$  and  $V(G)$  copies of  $H$  and joining the  $i$ th vertex of  $G$  to every vertex in the  $i$ th copy of  $H$  [3,7]. We will omit the subscript  $G$  when the graph is clear from the context.

A topological index is a numerical descriptor of the molecular structure derived from the corresponding (hydrogen-depleted) molecular graph. Various topological indices are widely used for quantitative structure-property relationship (QSPR) and quantitative structure-activity relationship (QSAR) studies [2,14,16-17,19].

The eccentric connectivity index of  $G$ ,  $\xi^c(G)$ , proposed by Sharma, Goswami and Madan [24], has been employed successfully for the development of numerous mathematical models for the prediction of biological activities of diverse nature [21-22]. It is defined as

$$\xi^c(G) = \sum_{v \in V(G)} \deg_G(v) ec_G(v),$$

where  $\deg_G(v)$  denotes the degree of the vertex  $v$  in  $G$  and  $ec_G(v)$  is the largest distance between  $v$  and any other vertex  $u$  of  $G$ . The quantity  $ec_G(v)$  is named as the eccentricity of vertex  $v$  in  $G$ . In fact, one can rewrite the eccentric connectivity index as

\*To whom all correspondence should be sent:

E-mail: zeynep.berberler@deu.edu.tr; murat.berberler@deu.edu.tr

$$\xi^c(G) = \sum_{uv \in E(G)} (ec(u) + ec(v)).$$

The total eccentricity of a graph  $G$  [18], denoted by  $\theta(G)$ , is the sum of all the eccentricities of  $G$ , that is,

$$\theta(G) = \sum_{v \in V(G)} ec(v).$$

The eccentric connectivity and total eccentricity polynomials of  $G$  are defined as

$$\Xi(G, x) = \sum_{v \in V(G)} \deg_G(v) x^{ec_G(v)}, \text{ and}$$

$$\Theta(G, x) = \sum_{v \in V(G)} x^{ec_G(v)},$$

respectively. It is easy to see that the eccentric connectivity index and the total eccentricity of a graph can be obtained from the corresponding polynomials by evaluating their first derivatives at  $x = 1$ .

The edge version of eccentric connectivity index of  $G$ ,  $\xi_e^c(G)$ , proposed by Xu and Guo [25], is defined as

$$\xi_e^c(G) = \sum_{f \in E(G)} \deg_G(f) \varepsilon_{c_G}(f),$$

where  $\deg_G(f)$  is the degree of an edge  $f$  and  $\varepsilon_{c_G}(f)$  is its eccentricity.

The Zagreb indices are the most known and widely used topological indices. The first and second Zagreb index of a graph were first introduced by Gutman and Trinajstić [13], defined as

$$M_1(G) = \sum_{v \in V(G)} \deg(v)^2,$$

$$M_2(G) = \sum_{uv \in E(G)} \deg(u) \deg(v),$$

respectively.

The first Zagreb index can be also expressed as a sum over edges of  $G$ ,

$$M_1(G) = \sum_{uv \in E(G)} (\deg_G(u) + \deg_G(v))$$

It is important and of interest to understand how certain invariants of composite graphs are related to the corresponding invariants of the original graphs. The concept of thorny graphs was proposed by Gutman [12]. Let  $G$  be a connected graph on  $n$  vertices. The thorn graph  $G^*$  of  $G$  is obtained by attaching  $p_i$  new vertices of degree one to the vertex  $u_i$  of the graph  $G$ , where  $p_i > 0$  and  $i = 1, \dots, n$ . The  $p_i$  pendent vertices attached to

vertex  $u_i$  are called the thorns of  $u_i$ . Recall the definition of the corona product. In particular, if  $p_1 = p_2 = \dots = p_n = p$ , then the thorny graph  $G^* = G \circ \overline{K_p}$ , where  $\overline{K_p}$  denotes the complement of a complete graph  $K_p$ .

The concept of thorn graphs was found in a variety of chemical applications as given in Refs. [4-6,11-12,15,20]. Interesting classes of graphs can also be obtained by specializing the parent graph in thorny graphs. Plerograms and kenograms are the two types of molecular graphs. Molecular graphs will be depicted in a usual way – atoms will be replaced by vertices and bonds by edges. In modern chemical graph theory, the plerograms are molecular graphs in which all atoms are represented by vertices whilst the kenograms are referred to as a hydrogen-suppressed or hydrogen-depleted molecular graphs [1]. The name thorn graphs for plerograms was introduced by Gutman [12]. What Cayley [1] calls a plerogram and Pólya [10] a C-H graph is just a thorn graph [12] and the parent graph would then be referred to as kenogram or a C-graph. In addition, caterpillars are the thorn graphs whose parent graph is a path [9]. For a given graph  $G$ , its  $t$ -fold bristled graph  $Brs_t(G)$  is obtained by attaching  $t$  vertices of degree 1 to each vertex of  $G$ . This graph can be represented as the corona of  $G$  and the complement graph on  $t$  vertices  $\overline{K_t}$ . The  $t$ -fold bristled graph of a given graph is also known as its  $t$ -thorny graph.

The paper proceeds as follows. In Section 2, the edge eccentric connectivity index of thorn graphs is investigated. Then, the results are applied to compute this eccentricity-related invariant for some important classes of graphs in terms of the underlying parent graph by specializing components and consider some special cases. A polynomial time algorithm is proposed in order to calculate the edge eccentric connectivity index for any simple finite connected graph without loops and multiple edges in Section 3.

## MAIN RESULTS

Let  $G$  be an  $(n, m)$ -graph and  $G^*$  be the thorn graph of  $G$  with parameters  $p_i$ ,  $p_i > 0$ ,  $i = 1, \dots, n$ . We let  $t_{ij}$  be the  $j$ th thorn attached to the vertex  $u_i$  of  $G$ , where  $j = 1, \dots, p_i$ . By the definition of thorn graphs, for every edge  $f$  in  $G^*$ , it holds that

$$\deg_{G^*}(f) = \begin{cases} \deg_G(u_i) + p_i + \deg_G(u_j) + p_j - 2, & \text{if } f = u_i u_j \in E(G); \\ \deg_G(u_i) + p_i - 1, & \text{if } f = u_i t_{ij} \text{ is a pendent edge in } G^*. \end{cases}$$

and

$$ec_{G^*}(f) = \begin{cases} ec_G(f) + 1, & \text{if } f \in E(G); \\ ec_G(u_i), & \text{if } f = u_i t_{ij} \text{ is a pendent edge in } G^*. \end{cases}$$

Now, we start to compute the edge version of eccentric connectivity index of thorn graphs.

*Theorem 2.1.* Let  $G$  be an  $(n, m)$ -graph. If  $G^*$  is the thorn graph of  $G$  with parameters  $p_i, p_i > 0, i = 1, \dots, n$ , then

$$\begin{aligned} \xi_e^c(G^*) &= \xi_e^c(G) + M_1(G) - 2m + \\ &\sum_{f=u_i u_j \in E(G)} (p_i + p_j)(ec_G(u_i u_j) + 1) + \\ &\sum_{i=1}^n p_i (\deg_G(u_i) + p_i - 1) ec_G(u_i) \end{aligned}$$

*Proof.* By the definition of edge eccentric connectivity index, we have:

$$\begin{aligned} \xi_e^c(G^*) &= \sum_{f \in E(G^*)} \deg_{G^*}(f) ec_{G^*}(f) \\ &= \sum_{f=u_i u_j \in E(G)} \deg_{G^*}(u_i u_j) ec_{G^*}(u_i u_j) + \\ &\sum_{i=1}^n \sum_{j=1}^{p_i} \deg_{G^*}(u_i t_{ij}) ec_{G^*}(u_i t_{ij}) \end{aligned}$$

By substituting the values of parameters due to the definition of thorn graphs, we compute

$$\begin{aligned} \xi_e^c(G^*) &= \sum_{f=u_i u_j \in E(G)} (\deg_G(u_i) + p_i + \deg_G(u_j) + p_j - 2)(ec_G(u_i u_j) + 1) + \\ &\sum_{i=1}^n \sum_{j=1}^{p_i} (\deg_G(u_i) + p_i - 1) ec_G(u_i) \\ &= \sum_{f=u_i u_j \in E(G)} (\deg_G(u_i) + \deg_G(u_j) - 2) ec_G(u_i u_j) + \\ &\sum_{f=u_i u_j \in E(G)} (\deg_G(u_i) + \deg_G(u_j) - 2) + \end{aligned}$$

$$\begin{aligned} &\sum_{f=u_i u_j \in E(G)} (p_i + p_j) ec_G(u_i u_j) + \sum_{f=u_i u_j \in E(G)} (p_i + p_j) + \\ &\sum_{i=1}^n p_i (\deg_G(u_i) + p_i - 1) ec_G(u_i) \end{aligned}$$

Since

$$\sum_{e \in E(G)} \deg_G(e) = \sum_{v \in V(G)} (\deg_G(v))^2 - 2|E(G)| \quad [8],$$

we get

$$\begin{aligned} \xi_e^c(G^*) &= \xi_e^c(G) + M_1(G) - 2m + \\ &\sum_{f=u_i u_j \in E(G)} (p_i + p_j)(ec_G(u_i u_j) + 1) + \\ &\sum_{i=1}^n p_i (\deg_G(u_i) + p_i - 1) ec_G(u_i) \end{aligned}$$

Hence, the desired result of  $\xi_e^c(G^*)$  holds.

The following corollaries for thorny graphs follow easily by direct calculations.

*Corollary 2.1.* If  $G^*$  is the thorn graph of an  $(n, m)$ -graph  $G$  with parameters  $p_1 = p_2 = \dots = p_n = p$ , then

$$\begin{aligned} \xi_e^c(G^*) &= \xi_e^c(G) + M_1(G) + 2m(p-1) + \\ &p(\xi_e^c(G) + (p-1)\theta(G)) + \\ &2p \sum_{f=u_i u_j \in E(G)} ec_G(u_i u_j) \end{aligned}$$

*Corollary 2.2.* If  $G^*$  is the thorn graph of an  $(n, m)$ -graph  $G$  with parameters  $p_1 = p_2 = \dots = p_n = 1$ , then

$$\begin{aligned} \xi_e^c(G^*) &= \xi_e^c(G) + M_1(G) + \\ &\xi_e^c(G) + 2 \sum_{f=u_i u_j \in E(G)} ec_G(u_i u_j) \end{aligned}$$

*Observation 2.1.*

By Corollary 2.1, the edge eccentric connectivity index of  $t$ -fold bristled graph of a given graph  $G$  can easily be computed.

*Example 2.1.* Let  $G$  be an  $(n, m)$ -graph. Then

$$\begin{aligned} \xi_e^c(Brs_t(G)) &= \xi_e^c(G) + M_1(G) + 2m(t-1) + \\ &t(\xi_e^c(G) + (t-1)\theta(G)) + \\ &2t \sum_{f=u_i u_j \in E(G)} ec_G(u_i u_j) \end{aligned}$$

For the simple bristled graph ( $t = 1$ )

$$\xi_e^c(Brs_t(G)) = \xi_e^c(G) + M_1(G) + \sum_{f=u_i u_j \in E(G)} ec_G(u_i u_j)$$

$$\xi_e^c(Brs_t(C_n)) = \begin{cases} n(n(1+t(t+5)/2) - 2t), & \text{if } n \text{ is even;} \\ n(n-1+t(nt-t+3n-3)/2), & \text{if } n \text{ is odd.} \end{cases}$$

From the above formula and Observation 2.1, the edge eccentric connectivity index of the  $t$ -fold bristled graph of  $P_n$  and  $C_n$  can easily be computed.

Note that the 2-fold bristled graph of  $P_4$  is the molecular graph related to polyethene.

$$\xi_e^c(Brs_t(P_n)) = \begin{cases} \frac{3n(n(t^2+3t+2)-8) - t(2n(t+13)+t-17)+26}{4}, & \text{if } n \text{ is odd;} \\ \frac{n(3n(t^2+3t+2) - 2(t^2+13t+12))}{4+4t+6}, & \text{if } n \text{ is even.} \end{cases}$$

and

**Table 2.1.** Invariants of path and cycle

$G$	$P_n$	$C_n$
$\xi_e^c(G)$	$\begin{cases} (3n^2 - 16n + 21)/2, \\ \text{if } n \text{ is odd;} \\ (3n^2 - 16n + 20)/2, \\ \text{if } n \text{ is even.} \end{cases}$	$2n \lfloor (n-2)/2 \rfloor$
$\xi^c(G)$	$\begin{cases} (3n^2 - 6n + 4)/2, \\ \text{if } n \text{ is even;} \\ 3(n-1)^2/2, \\ \text{if } n \text{ is odd.} \end{cases}$	$\begin{cases} n^2, \\ \text{if } n \text{ is even;} \\ n(n-1), \\ \text{if } n \text{ is odd.} \end{cases}$
$\theta(G)$	$\begin{cases} 3/4 n^2 - 1/2 n, \\ \text{if } n \text{ is even;} \\ 3/4 n^2 - 1/2 n - 1/4, \\ \text{if } n \text{ is odd.} \end{cases}$	$\begin{cases} 1/2 n^2, \\ \text{if } n \text{ is even;} \\ 1/2 n(n-1), \\ \text{if } n \text{ is odd.} \end{cases}$
$M_1(G)$	$4n - 6$	$4n$
$\sum_{f=u_i u_j \in E(G)} ec_G(u_i u_j)$	$\begin{cases} (3n^2 - 12n + 9)/4, \\ \text{if } n \text{ is odd;} \\ (3n^2 - 12n + 8)/4, \\ \text{if } n \text{ is even.} \end{cases}$	$n \lfloor (n-2)/2 \rfloor$

COMPUTING THE EDGE ECCENTRIC CONNECTIVITY INDEX OF A GRAPH

In this section, a polynomial time algorithm is proposed in order to calculate the edge eccentric

connectivity index for any simple finite undirected connected graph without loops and multiple edges by using the exploration algorithm *BFS* [23]. The adjacency matrix  $A$  is used to store the neighbors of each vertex.

**Algorithm.** The Edge Eccentric Connectivity Index

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```

for all  $v \in V$  do
     $BFS(G, v)$ ;
     $\delta[v] \leftarrow 0$ ;
    for all  $u \in V$  do
         $D[v, u] \leftarrow dist[u]$ ;
         $\delta[v] \leftarrow \delta[v] + A[v, u]$ ;
    end
end
 $\xi_e^c(G) \leftarrow 0$ ;
for all  $f_1(u_1, v_1) \in E$  do
     $deg_G(f_1) \leftarrow \delta[u_1] + \delta[v_1] - 2$ 
    for all  $f_2(u_2, v_2) \in E$  do
         $d_G[f_2] \leftarrow \min\{D[u_1, u_2], D[u_1, v_2], D[v_1, u_2], D[v_1, v_2]\}$ 
        end
         $ec_G(f_1) \leftarrow \max\{d_G[f] \mid \forall f \in E\}$ 
         $\xi_e^c(G) \leftarrow \xi_e^c(G) + deg_G(f_1) * ec_G(f_1)$ 
    end
end

```

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The below function  $BFS$  returns the distances from a source vertex  $v$  to all other vertices in the graph.

---

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function  $BFS(G, v)$ ;

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**Input:** Graph  $G = (V, E)$ , start vertex  $v$

**Output:** for all vertices  $u$  reachable from  $v$ ,  $dist[u]$  is set to the distance from  $v$  to  $u$

```

for all  $u \in V$  do
     $dist[u] \leftarrow \infty$ ;
end
 $dist[v] \leftarrow 0$ ;
 $Q \leftarrow [v]$ ;
while  $Q$  is not empty do
     $u \leftarrow eject(Q)$ ;
    for all edges  $(u, w) \in E$  do
        if  $dist[v] = \infty$  then
             $INJECT(Q, v)$ ;
             $dist[w] \leftarrow dist[u] + 1$ ;
        end
    end
end

```

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The Edge Eccentric Connectivity Index calculation algorithm runs in polynomial time. The running time of function  $BFS(G, v)$  is  $O(|V| + |E|)$  and the function  $BFS(G, v)$  runs for all  $v \in V(G)$ . Since the binary combinations of the edge set, that are  $C(|E|, 2) = |E|(|E| - 1)/2$ , are handled one by one in the algorithm and since  $|E| \leq |V|(|V| - 1)/2$ , the complexity for the algorithm described is  $O(|E|^2) = O(|V|^4)$

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## Индекс на ексцентрична реброва свързаност на наноторни

З. Н. Берберлер, М. Е. Берберлер

Научен факултет, Департамент по компютърни науки, Университет Докуз Ейлул, 65160, Измир, Турция

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Индексът на ексцентрична реброва свързаност на един граф  $G$  се дефинира като  $\xi_e^c(G) = \sum_{f \in E(G)} \deg_G(f) \varepsilon_{c_G}(f)$ , където  $\deg_G(f)$  е степента на едно ребро  $f$  и  $\varepsilon_{c_G}(f)$  е неговата ексцентричност. В тази статия изследваме индекса на ексцентрична реброва свързаност на един клас нанографи, а именно наноторни и представяме една експлицитна формула за индекса на ексцентрична реброва свързаност на наноторни. Резултатите се прилагат за изчисляване на този инвариант, свързан с ексцентричността за някои важни класове химически графи от специализирани компоненти на този тип графи. В допълнение, е конструиран алгоритъм, предназначен за изчисляване на индекса на ексцентричната реброва свързаност на графи.